

## CS 1155 Spring 1999 HW#7 Sample Solutions

1. Consider the following recurrence equation:  $s_n = 6s_{n-1} - 11s_{n-2} + 6s_{n-3}$ , with  $s_0 = 3$ ,  $s_1 = 6$ , and  $s_2 = 14$ .

(a) (5 points) What is the characteristic equation of this recurrence?

Solution:

It is reasonable to hope that  $s^n$  would have the form  $s_n = cr^n$  for a nonzero constant  $c$ . This would mean that

$$cr^n = 6cr^{n-1} - 11cr^{n-2} + 6cr^{n-3}$$

Dividing by  $cr^{n-3}$  would then give

$$r^3 = 6r^2 - 11r + 6$$

Therefore, the characteristic equation of the recurrence relation would be

$$x^3 - 6x^2 + 11x - 6 = 0$$

(b) (5 points) What are the roots of the characteristic equation?

Solution:

The characteristic equation is equivalent to the following:

$$r^3 - 6r^2 + 9r - 6 = 0$$

$$(r - 3)^2 r + 2(r - 3) = 0$$

$$(r - 3)[(r - 3)r + 2] = 0$$

$$(r - 3)(r - 2)(r - 1) = 0$$

Clearly the roots of the characteristic equation are 1, 2, and 3.

(c) (10 points) What is the solution of the recurrence? Prove your answer.

Solution:

Since the characteristic equation has three different roots, we can assume that

$$s_n = c_1 r_1^n + c_2 r_2^n + c_3 r_3^n$$

We know that  $s_0 = 3$ ,  $s_1 = 6$ , and  $s_2 = 14$ . Let  $r_1 = 1$ ,  $r_2 = 2$ , and  $r_3 = 3$ . Solving

$$\begin{cases} 3 = c_1 + c_2 + c_3 \\ 6 = c_1 + 2c_2 + 3c_3 \\ 14 = c_1 + 4c_2 + 9c_3 \end{cases}$$

gives us  $c_1 = c_2 = c_3 = 1$ . In other words, the solution of the recurrence is

$$s_n = r_1^n + r_2^n + r_3^n = 1 + 2^n + 3^n$$

We can prove this by using induction. The basis cases where  $n = 0, 1, 2$ , can be verified easily. Assuming that for any  $n$  so that  $0 \leq n \leq m$ , our solution is correct, we want to prove that for  $n = m + 1$ , our solution is also correct:

$$\begin{aligned} s_n &= s_{m+1} \\ &= 6s_m - 11s_{m-1} + 6s_{m-2} \\ &= 6(1 + 2^m + 3^m) - 11(1 + 2^{m-1} + 3^{m-1}) + 6(1 + 2^{m-2} + 3^{m-2}) \\ &= (6 - 11 + 6) + (6 \times 2^2 - 11 \times 2 + 6) \times 2^{m-2} + (6 \times 3^2 - 11 \times 3 + 6) \times 3^{m-2} \end{aligned}$$

$$= 1 + 2^{m+1} + 3^{m+1}$$

Therefore, our solution of the recurrence is correct. ■

2. (10 points) Consider the sequence defined in Example 4, page 239 of the text. The characteristic equation has only one root, namely  $r = 3$ . (Therefore, for the solution of this recurrence, part (b) of Theorem 1 is applicable.) The proof of part (a) of Theorem 1 has two steps: (a) a basis step where the constants  $c_1, c_2$  are determined, and (b) an inductive step. Which of these two steps in the proof of part (a) of Theorem 1 fails for this example? Why?

Solution:

The first step fails because the equations

$$\begin{cases} s_0 = c_1 + c_2 \\ s_1 = c_1 r_1 + c_2 r_2 \end{cases}$$

cannot be solved for  $c_1$  or  $c_2$  if the characteristic equation has only one root, i.e.  $r_1 = r_2$ . ■

3. (15 points) Consider the matrices  $M^1, M^2$ , and  $M^n$  defined on page 231 of the text book. Using these definitions, prove that, for all  $n \geq 1$ ,  $FIB(n+1) \cdot FIB(n-1) - FIB^2(n) = (-1)^n$ .

Solution:

We could use induction for this problem, but a simpler route exists:

$$\begin{aligned} & FIB(n+1) \cdot FIB(n-1) - FIB^2(n) \\ &= \left| \begin{bmatrix} FIB(n+1) & FIB(n) \\ FIB(n) & FIB(n-1) \end{bmatrix} \right| = |M^n| = |M^1|^n = \left| \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix} \right|^n = (-1)^n \end{aligned}$$

4. (10 points) Consider the sequence  $s_n$  where  $s_0 = 2, s_1 = 1$  and  $s_n = s_{n-1} + s_{n-2}$  for  $n \geq 2$ .  
(a) Calculate  $s_n$  for  $n = 2, 3, 4, 5$  and  $6$ .

Solution:

$$s_2 = s_1 + s_0 = 1 + 2 = 3$$

$$s_3 = s_2 + s_1 = 3 + 1 = 4$$

$$s_4 = s_3 + s_2 = 4 + 3 = 7$$

$$s_5 = s_4 + s_3 = 7 + 4 = 11$$

$$s_6 = s_5 + s_4 = 11 + 7 = 18$$

- (b) Give an explicit formula for  $s_n$ .

Solution:

Theorem 1 is applicable in this case. The characteristic equation is

$$x^2 - x - 1 = 0$$

It has two roots:  $\frac{1 + \sqrt{5}}{2}$  and  $\frac{1 - \sqrt{5}}{2}$ . Solving

$$\begin{cases} 2 = c_1 + c_2 \\ 1 = \frac{1 + \sqrt{5}}{2}c_1 + \frac{1 - \sqrt{5}}{2}c_2 \end{cases}$$

gives us  $c_1 = c_2 = 1$ . Therefore,

$$s_n = \left(\frac{1 + \sqrt{5}}{2}\right)^n + \left(\frac{1 - \sqrt{5}}{2}\right)^n$$

5. (10 points) Give an explicit formula for  $s_n$ .

(c)  $s_0 = 5$  and  $s_n = s_{n-1}$  for  $n \geq 1$ .

Solution:

$$s_n = 5$$

(d)  $s_0 = 3$ ,  $s_1 = 5$  and  $s_n = 2s_{n-1} - s_{n-2}$  for  $n \geq 2$ .

Solution:

Using Theorem 1, the characteristic equation is

$$x^2 - 2x + 1 = 0$$

which has only one root, 1. Solving

$$\begin{cases} 3 = c_1 + c_2 \cdot 0 \\ 5 = c_1 + c_2 \cdot 1 \end{cases}$$

we have  $c_1 = 3$  and  $c_2 = 2$ . Therefore

$$s_n = 3 + 2n$$

6. (10 points) Theorem 2 does not specifically apply to the following recurrence, but the ideas carry over.

(b) Given that  $t_{3n} = 3t_n + f(n)$ , find a formula for  $t_{3^m}$ .

Solution:

Our guess would be

$$s_{3^m} = 3^m \cdot \left( s_1 + \frac{1}{3} \sum_i^{m-1} \frac{f(3^i)}{3^i} \right) \tag{\$}$$

and we can prove it by using induction. The basis is

$$s_1 = s_{3^0} = 3^0 \cdot \left( s_1 + \frac{1}{3} \times 0 \right) = s_1$$

Assume that (§) is true for a specific  $m$ ,

$$\begin{aligned}
s_{3^{m+1}} &= s_3 \cdot 3^m \\
&= 3 \cdot s_{3^m} + f(3^m) \\
&= 3^{m+1} \cdot \left( s_1 + \frac{1}{3} \sum_i^{m-1} \frac{f(3^i)}{3^i} \right) + f(3^m) \\
&= 3^{m+1} \cdot \left( s_1 + \frac{1}{3} \sum_i^{m-1} \frac{f(3^i)}{3^i} + \frac{f(3^m)}{3 \cdot 3^m} \right) \\
&= 3^{m+1} \cdot \left( s_1 + \frac{1}{3} \sum_i^m \frac{f(3^i)}{3^i} \right)
\end{aligned}$$

The inductive step is also true. Therefore (§) is true for all  $m \geq 0$ . ■

7. (10 points) Recursively define  $a_0 = 1$ ,  $a_1 = 3$ ,  $a_2 = 5$  and  $a_n = 3a_{n-2} + 2a_{n-3}$  for  $n \geq 3$ .

(d) Prove that  $a_n = 2a_{n-1} + (-1)^{n-1}$  for  $n \geq 1$ .  (§)

Solution:

We need to use the *Second Principle of Mathematical Induction* for this problem. First, we need to verify that (§) is true for  $n = 1$  and  $2$ .

$$\text{if } n = 1, 2a_{n-1} + (-1)^{n-1} = 2a_{1-1} + (-1)^{1-1} = 2a_0 + 1 = 3 = a_1$$

$$\text{if } n = 2, 2a_{n-1} + (-1)^{n-1} = 2a_{2-1} + (-1)^{2-1} = 2a_1 - 1 = 5 = a_2$$

Then, assuming that (§) is true for indices from  $1$  to  $n - 1$ , we need to prove that (§) is still true for index  $n$ .

$$\begin{aligned}
a_n &= 3a_{n-2} + 2a_{n-3} \\
&= 3(2a_{n-3} + (-1)^{n-3}) + 2(2a_{n-4} + (-1)^{n-4}) \\
&= 6a_{n-3} + 3(-1)^{n-3} + 4a_{n-4} + 2(-1)^{n-4} \\
&= 6a_{n-3} + 4a_{n-4} + (-3 + 2)(-1)^{n-4} \\
&= 6a_{n-3} + 4a_{n-4} + (-1)^{n-3} \\
&= 2(3a_{n-3} + 2a_{n-4}) + (-1)^2(-1)^{n-3} \\
&= 2a_{n-1} + (-1)^{n-1}
\end{aligned}$$

The proof is not complete at this point. Notice that in order to use the (§) assumption in the second sub-step above,  $n - 3$  has to be  $\geq 1$ . This means that our inductive step works only if  $n \geq 4$  and the  $n = 3$  case has to be covered in the induction basis:

$$a_3 = 3a_1 + 2a_0 = 9 + 2 = 11$$

$$\text{if } n = 3, 2a_{n-1} + (-1)^{n-1} = 2a_{3-1} + (-1)^{3-1} = 2a_2 + 1 = 11 = a_3$$

Now the proof is complete. ■

8. Let  $p, q$  be two propositions. Consider the following recursive definition of the set  $E$  of well-formed conjunctive and disjunctive Boolean expressions using  $p$  and  $q$ :

**Basis:**  $p, q \in E$ .

**Recursive Step:** If  $e_1, e_2$  are expressions in  $E$ , then  $(e_1 \vee e_2) \in E$  and  $(e_1 \wedge e_2) \in E$ .

An expression is in  $E$  iff it is obtained from the basis by a finite number of applications of the recursive step.

- (a) (5 points) What are the expressions obtained after two applications of the recursive step to the set of expressions defined by the basis?

Solution:

Let  $E_i$  denote the set  $E$  after  $i$  applications of the recursive step.

$i$	$E_i$
0	$p, q$
1	$p, q,$ $(p \wedge p), (p \vee p), (q \wedge q), (q \vee q), (p \wedge q), (p \vee q), (q \wedge p), (q \vee p)$
2	$p, q,$ $(p \wedge p), (p \vee p), (q \wedge q), (q \vee q), (p \wedge q), (p \vee q), (q \wedge p), (q \vee p),$ plus all $10 \times 10 \times 2 = 200$ cases of $(e_1 \wedge e_2)$ and $(e_1 \vee e_2)$ , where $e_1, e_2 \in E_1$

- (b) (10 points) For an expression  $e$ , let  $P(e)$  denote the number of occurrences of the propositions  $p$  and  $q$  in  $e$  and let  $O(e)$  denote the number of occurrences of the operators  $\wedge$  and  $\vee$  in  $e$ . Prove, by induction, that for every expression  $e \in E$ ,  $P(e) = 1 + O(e)$ .

Solution:

Induction basis: for every expression  $e \in E_0 = \{p, q\}$ ,  $P(e) = 1 + O(e)$ , because

$$P(p) = 1 = 1 + O(p), P(q) = 1 = 1 + O(q)$$

Inductive step: assuming that for every expression  $e \in E_i$ ,  $P(e) = 1 + O(e)$ , where  $0 \leq i \leq n$ , we want to prove that for every expression  $e \in E_{n+1}$ ,

$$P(e) = 1 + O(e) \tag{\$}$$

If  $e \in E_n$ , (§) is obviously true according to the inductive assumption. Therefore we only need to consider the cases where  $e \in E_{n+1}$  but  $e \notin E_n$ . Then  $e$  can be either  $(e_1 \wedge e_2)$  or  $(e_1 \vee e_2)$ , where  $e_1 \in E_i, e_2 \in E_j, 0 \leq i \leq n, 0 \leq j \leq n$ . If  $e = (e_1 \wedge e_2)$ ,

$$\begin{aligned} P(e) &= P((e_1 \wedge e_2)) = P(e_1) + P(e_2) \\ &= (1 + O(e_1)) + (1 + O(e_2)) \\ &= 1 + (1 + O(e_1) + O(e_2)) \\ &= 1 + O((e_1 \wedge e_2)) \\ &= 1 + O(e) \end{aligned}$$

The  $e = (e_1 \vee e_2)$  case is very similar.

In summary, according to the *Second Principle of Mathematical Induction*, (§) is true for all  $E$ . ■