

CS 6260

Some number theory

Let $\mathbf{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$ denote the set of integers.
Let $\mathbf{Z}^+ = \{1, 2, \dots\}$ denote the set of positive integers and
 $\mathbf{N} = \{0, 1, 2, \dots\}$ the set of non-negative integers.

If a, N are integers with $N > 0$ then there are unique integers r, q such that $a = Nq + r$ and $0 \leq r < N$.

We associate to any positive integer N the following two sets:
 $\mathbf{Z}_N = \{0, 1, \dots, N-1\}$, $\mathbf{Z}_N^* = \{i \in \mathbf{Z} : 1 \leq i \leq N-1 \text{ and } \gcd(i, N) = 1\}$

Groups

- Def. Let G be a non-empty set and let \cdot denote a binary operation on G . We say that G is a group if it has the following properties:
 1. Closure: For every $a, b \in G$ it is the case that $a \cdot b$ is also in G .
 2. Associativity: For every $a, b, c \in G$ it is the case that $(a \cdot b) \cdot c = a \cdot (b \cdot c)$.
 3. Identity: There exists an element $1 \in G$ such that $a \cdot 1 = 1 \cdot a = a$ for all $a \in G$.
 4. Invertibility: For every $a \in G$ there exists a unique $b \in G$ such that $a \cdot b = b \cdot a = 1$.

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inverse, denoted a^{-1}

- Fact. Let N be a positive integer. Then \mathbf{Z}_N is a group under addition modulo N , and \mathbf{Z}_N^* is a group under multiplication modulo N .
- In any group, we can define an exponentiation operation:
 - if $i = 0$ then a^i is defined to be 1 ,
 - if $i > 0$ then $a^i = a \cdot a \cdot \dots \cdot a$ (i times)
 - if $i < 0$ then $a^i = a^{-1} \cdot a^{-1} \cdot \dots \cdot a^{-1}$ ($j = -i$ times)
- For all $a \in \mathbf{G}$ and all $i, j \in \mathbf{Z}$:
 - $a^{i+j} = a^i \cdot a^j$
 - $(a^i)^j = a^{ij}$
 - $a^{-i} = (a^i)^{-1} = (a^{-1})^i$

- The order of a group is its size
- Fact. Let \mathbf{G} be a group and let $m = |\mathbf{G}|$ be its order. Then $a^m = 1$ for all $a \in \mathbf{G}$
- Fact. Let \mathbf{G} be a group and let $m = |\mathbf{G}|$ be its order. Then $a^i = a^{i \bmod m}$ for all $a \in \mathbf{G}$ and all $i \in \mathbf{Z}$.
- Example. Let us work in the group $\mathbf{Z}_{21}^* = \{1, 2, 4, 5, 8, 10, 11, 13, 16, 17, 19, 20\}$ under the operation of multiplication modulo 21. $m=12$.

$$5^{86} \bmod 21 = 5^{86 \bmod 12} \bmod 21 = 5^2 \bmod 21 = 25 \bmod 21 = 4$$

- If \mathbf{G} is a group, a set $\mathbf{S} \subseteq \mathbf{G}$ is called a subgroup if it is a group in its own right, under the same operation as that under which \mathbf{G} is a group.
- If we already know that \mathbf{G} is a group, there is a simple way to test whether \mathbf{S} is a subgroup:
 - it is one if and only if $x \cdot y^{-1} \in \mathbf{S}$ for all $x, y \in \mathbf{S}$. Here y^{-1} is the inverse of y in \mathbf{G} .
- Fact. Let \mathbf{G} be a group and let \mathbf{S} be a subgroup of \mathbf{G} . Then the order of \mathbf{S} divides the order of \mathbf{G} .

Algorithms and their running times

- Since in cryptography we will be working with BIG numbers, the complexity of algorithms taking numbers as inputs is measured as a function of the bit-length of the numbers.
- E.g. PrintinBinary (A), where $A=2^k$ takes k operations

Some basic algorithms

Algorithm	Input	Output	Running Time
INT-DIV	a, N ($N > 0$)	(q, r) with $a = Nq + r$ and $0 \leq r < N$	$O(a \cdot N)$
MOD	a, N ($N > 0$)	$a \bmod N$	$O(a \cdot N)$
EXT-GCD	a, b ($(a, b) \neq (0, 0)$)	(d, \bar{a}, \bar{b}) with $d = \gcd(a, b) = a\bar{a} + b\bar{b}$	$O(a \cdot b)$
MOD-ADD	a, b, N ($a, b \in \mathbf{Z}_N$)	$(a + b) \bmod N$	$O(N)$
MOD-MULT	a, b, N ($a, b \in \mathbf{Z}_N$)	$ab \bmod N$	$O(N ^2)$
MOD-INV	a, N ($a \in \mathbf{Z}_N^*$)	$b \in \mathbf{Z}_N^*$ with $ab \equiv 1 \pmod{N}$	$O(N ^2)$
MOD-EXP	a, n, N ($a \in \mathbf{Z}_N$)	$a^n \bmod N$	$O(n \cdot N ^2)$
EXP $_G$	a, n ($a \in G$)	$a^n \in G$	$2 n $ G -operations

Cyclic groups and generators

- If $g \in \mathbf{G}$ is any member of the group, the order of g is defined to be the least positive integer n such that $g^n = 1$.
We let $\langle g \rangle = \{ g^i : i \in \mathbf{Z}_n \} = \{ g^0, g^1, \dots, g^{n-1} \}$ denote the set of group elements generated by g . This is a subgroup of order n .
- Def.** An element g of the group is called a generator of \mathbf{G} if $\langle g \rangle = \mathbf{G}$, or, equivalently, if its order is $m = |\mathbf{G}|$.
- Def.** A group is cyclic if it contains a generator.
- If g is a generator of \mathbf{G} , then for every $a \in \mathbf{G}$ there is a unique integer $i \in \mathbf{Z}_m$ such that $g^i = a$. This i is called the discrete logarithm of a to base g , and we denote it by $\text{DLog}_{\mathbf{G},g}(a)$.
- $\text{DLog}_{\mathbf{G},g}(a)$ is a function that maps \mathbf{G} to \mathbf{Z}_m , and moreover this function is a bijection.
- The function of \mathbf{Z}_m to \mathbf{G} defined by $i \rightarrow g^i$ is called the discrete exponentiation function

- Example.** Let $p = 11$. Then $\mathbf{Z}_{11}^* = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$ has order $p - 1 = 10$. We find the subgroups generated by group elements 2 and 5. We raise them to the powers $0, \dots, 9$.

i	0	1	2	3	4	5	6	7	8	9
$2^i \text{ mod } 11$	1	2	4	8	5	10	9	7	3	6
$5^i \text{ mod } 11$	1	5	3	4	9	1	5	3	4	9

$$\langle 2 \rangle = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\} = \mathbf{Z}_{11}^* \quad \langle 5 \rangle = \{1, 3, 4, 5, 9\}$$

2 is a generator and thus \mathbf{Z}_{11}^* is cyclic.

a	1	2	3	4	5	6	7	8	9	10
$\text{DLog}_{\mathbf{Z}_{11}^*, 2}(a)$	0	1	8	2	4	9	7	3	6	5

Choosing cyclic group and generators

- The discrete log function is conjectured to be one-way (hard to compute) for some cyclic groups \mathbf{G} . Due to this fact we often seek cyclic groups.
- Examples of cyclic groups:
 - \mathbf{Z}_p^* for a prime p ,
 - a group of prime order
- We will also need generators. How to choose a candidate and test it?
- Fact.** Let \mathbf{G} be a cyclic group and let $m = |\mathbf{G}|$. Let $p_1^{\alpha_1} \cdot \dots \cdot p_n^{\alpha_n}$ be the prime factorization of m and let $m_i = m/p_i$ for $i = 1, \dots, n$.
Then $g \in \mathbf{G}$ is a generator of \mathbf{G} if and only if for all $i = 1, \dots, n$: $g^{m_i} \neq 1$.
- Fact.** Let \mathbf{G} be a cyclic group of order m , and let g be a generator of \mathbf{G} . Then $\text{Gen}(\mathbf{G}) = \{ g^i \in \mathbf{G} : i \in \mathbf{Z}_m^* \}$ and $|\text{Gen}(\mathbf{G})| = \phi(m)$.

- Example.** Let us determine all the generators of the group \mathbf{Z}_{11}^* . Its size is $m = \phi(11) = 10$, and the prime factorization of 10 is $2^1 \cdot 5^1$. Thus, the test for whether a given $a \in \mathbf{Z}_{11}^*$ is a generator is that $a^2 \neq 1 \pmod{11}$ and $a^5 \neq 1 \pmod{11}$.

a	1	2	3	4	5	6	7	8	9	10
$a^2 \text{ mod } 11$	1	4	9	5	3	3	5	9	4	1
$a^5 \text{ mod } 11$	1	10	1	1	1	10	10	10	1	10

- $\text{Gen}(\mathbf{Z}_{11}^*) = \{2, 6, 7, 8\}$.
- Double-checking: $|\mathbf{Z}_{11}^*| = 10$, $\mathbf{Z}_{10}^* = \{1, 3, 7, 9\}$
 $\{ 2^i \in \mathbf{G} : i \in \mathbf{Z}_{10}^* \} = \{ 2^1, 2^3, 2^7, 2^9 \pmod{11} \} = \{2, 6, 7, 8\}$

Algorithm for finding a generator

- The most common choice of a group in crypto is \mathbf{Z}_p^* for a prime p .
- Idea.** Pick a random element and test it. Chose p s.t. the prime factorization of the order of the group $(p-1)$ is known. E.g., chose a prime p s.t. $p=2q+1$ for some prime q .
- Algorithm FIND-GEN(p)
 - $q \leftarrow (p-1)/2$
 - found $\leftarrow 0$
 - While (found $\neq 1$) do
 - $g \leftarrow \mathbf{Z}_p^* - \{1, p-1\}$
 - If $(g^2 \bmod p \neq 1)$ and $(g^q \bmod p \neq 1)$ then found $\leftarrow 1$
 - EndWhile
 - Return g
- The probability that an iteration of the algorithm is successful in finding a generator is

$$\frac{|\text{Gen}(\mathbf{Z}_p^*)|}{|\mathbf{Z}_p^*| - 2} = \frac{\varphi(p-1)}{p-3} = \frac{\varphi(2q)}{2q-2} = \frac{q-1}{2q-2} = \frac{1}{2}$$

Squares and non-squares

- Def.** An element a of a group \mathbf{G} is called a square, or quadratic residue if it has a square root, meaning there is some $b \in \mathbf{G}$ such that $b^2 = a$ in \mathbf{G} .
- We let $\text{QR}(\mathbf{G}) = \{g \in \mathbf{G} : g \text{ is quadratic residue in } \mathbf{G}\}$
- We are mostly interested in the case where the group \mathbf{G} is \mathbf{Z}_N^* for some integer N .
- Defs.** An integer a is called a square mod N or quadratic residue mod N if $a \bmod N$ is a member of $\text{QR}(\mathbf{Z}_N^*)$. If $b^2 = a \pmod{N}$ then b is called a square-root of $a \bmod N$. An integer a is called a non-square mod N or quadratic non-residue mod N if $a \bmod N$ is a member of $\mathbf{Z}_N^* - \text{QR}(\mathbf{Z}_N^*)$.

- Def.** Let p be a prime. Define the Legendre symbol of a

$$J_p(a) = \begin{cases} 1 & \text{if } a \text{ is a square mod } p \\ 0 & \text{if } a \bmod p = 0 \\ -1 & \text{otherwise.} \end{cases}$$

- Example.** $\text{QR}(\mathbf{Z}_{11}^*)$?

a	1	2	3	4	5	6	7	8	9	10
$a^2 \bmod 11$	1	4	9	5	3	3	5	9	4	1

$$\text{QR}(\mathbf{Z}_{11}^*) = \{1, 3, 4, 5, 9\}$$

Recall that \mathbf{Z}_{11}^* is cyclic and 2 is a generator.

Fact. A generator is always a non-square. (But not all non-squares are generators).

a	1	2	3	4	5	6	7	8	9	10
$\text{DLog}_{\mathbf{Z}_{11}^*, 2}(a)$	0	1	8	2	4	9	7	3	6	5
$J_{11}(a)$	1	-1	1	1	1	-1	-1	-1	1	-1

- Fact.** Let $p \geq 3$ be a prime and let g be a generator of \mathbf{Z}_p^* . Then $\text{QR}(\mathbf{Z}_p^*) = \{g^i : i \in \mathbf{Z}_{p-1} \text{ and } i \text{ is even}\}$, and $|\text{QR}(\mathbf{Z}_p^*)| = (p-1)/2$

Facts. Let $p \geq 3$ be a prime. Then

- $J_p(a) \equiv a^{\frac{p-1}{2}} \pmod{p}$ for any $a \in \mathbf{Z}_p^*$
- $g^{\frac{p-1}{2}} \equiv -1 \pmod{p}$ for any generator g of \mathbf{Z}_p^*
- $J_p(ab \bmod p) = J_p(a) \cdot J_p(b)$ for any $a \in \mathbf{Z}_p^*$
- $J_p(g^{xy} \bmod p) = 1$ if and only if $J_p(g^x \bmod p) = 1$ or $J_p(g^y \bmod p) = 1$ for any generator g of \mathbf{Z}_p^* and any $x, y \in \mathbf{Z}_{p-1}$
- $\Pr \left[x \leftarrow \mathbf{Z}_{p-1}; y \leftarrow \mathbf{Z}_{p-1} : J_p(g^{xy}) = 1 \right] = 3/4$ for any generator g of \mathbf{Z}_p^*

Groups of prime order

- Def. An element h of a group \mathbf{G} is called non-trivial if it is not equal to the identity element of the group.
- Fact. Any non-trivial member of a group of prime order is a generator of the group.
- Fact. Let $q \geq 3$ be a prime such that $p = 2q + 1$ is also prime. Then $\text{QR}(\mathbf{Z}_p^*)$ is a group of prime order q . Furthermore, if g is any generator of \mathbf{Z}_p^* , then $g^2 \bmod p$ is a generator of $\text{QR}(\mathbf{Z}_p^*)$.
- Fact. Let g be a generator of a group of prime order q . Then for any element Z of the group

$$\Pr \left[x \stackrel{\$}{\leftarrow} \mathbf{Z}_q ; y \stackrel{\$}{\leftarrow} \mathbf{Z}_q : g^{xy} = Z \right] = \begin{cases} \frac{1}{q} \left(1 - \frac{1}{q}\right) & \text{if } Z \neq 1 \\ \frac{1}{q} \left(2 - \frac{1}{q}\right) & \text{if } Z = 1 \end{cases}$$

- Example. Let $q = 5$ and $p = 2q + 1 = 11$.

- $\text{QR}(\mathbf{Z}_{11}^*) = \{1, 3, 4, 5, 9\}$

We know that 2 is a generator of \mathbf{Z}_{11}^*

Let's verify that $4 = 2^2$ is a generator of $\text{QR}(\mathbf{Z}_{11}^*)$.

i	0	1	2	3	4
$4^i \bmod 11$	1	4	5	9	3