Sorting + Aggregation

Z

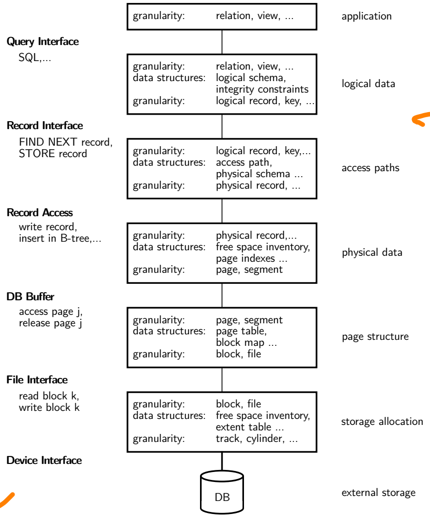
- relative
- finde Brüder

- ≥ 1B große
- Gruppierungs

Recap

A More Detailed Architecture

Query Exec



granularity: relation, view, ...

application

Query Interface

SQL,...

granularity: relation, view, ...
data structures: logical schema, integrity constraints
granularity: logical record, key, ...

logical data

Record Interface

FIND NEXT record, STORE record

granularity: logical record, key,...
data structures: access path, physical schema ...
granularity: physical record, ...

access paths

Record Access

write record, insert in B-tree,...

granularity: physical record,...
data structures: free space inventory, page indexes ...
granularity: page, segment

physical data

DB Buffer

access page j, release page j

granularity: page, segment
data structures: page table, block map ...
granularity: block, file

page structure

File Interface

read block k, write block k

granularity: block, file
data structures: free space inventory, extent table ...
granularity: track, cylinder, ...

storage allocation

Device Interface

DB

external storage

Indexing

Storage + Buffer Management

Anatomy of a Database System [Monologue]

- Process Manager
  - Connection Manager + Admission Control
- Query Processor
  - Query Parser
  - Query Optimizer (*a.k.a.*, Query Planner)
  - Query Executor
- Transactional Storage Manager
  - Lock Manager
  - Access Methods (*a.k.a.*, Indexes)
  - Buffer Pool Manager
  - Log Manager
- Shared Utilities
  - Memory, Disk, and Networking Manager

Query Processing

CS

8803

Query Execution

- We are now going to talk about how to execute queries using table heaps and indexes.
- Coming weeks:
    - Operator Algorithms
    - Query Processing Models
    - Runtime Architectures

Proj. Assignment

Iterator model

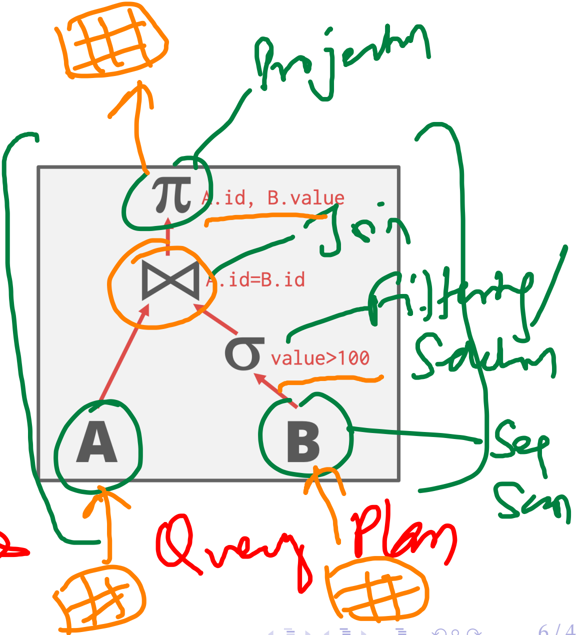
Tuple-at-a-time model

Query Plan

- The operators are arranged in a tree.
- Data flows from the leaves of the tree up towards the root.
- The output of the root node is the result of the query.

```
SELECT A.id, B.value
  FROM A, B
  WHERE A.id = B.id AND B.value > 100
```

Query →



Disk-Oriented DBMS

Buffer Manager
Access Methods

- We **cannot** assume that the results of a query fits in memory.
- We are going use the **buffer pool** to implement query execution algorithms that need to spill to disk.
- We are also going to prefer algorithms that maximize the amount of **sequential access**.

Today's Agenda

- External Merge Sort
- Tree-based Sorting
- Aggregation

Sorting

Aggregation

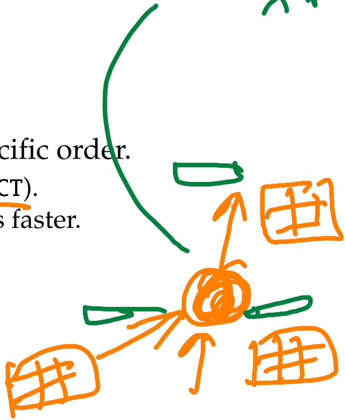
Join

External Merge Sort

Why do we need to sort?

tuple-at-a-time.

- Tuples in a table have no specific order.
- But queries often want to retrieve tuples in a specific order.
    - Trivial to support duplicate elimination (`DISTINCT`).
    - Bulk loading sorted tuples into a B+Tree index is faster.
    - Aggregation (`GROUP BY`).



Sorting Algorithms

- If data fits in memory, then we can use a standard in-memory sorting algorithm like **quick-sort**.
- If data does not fit in memory, then we need to use a technique that is aware of the cost of writing data out to disk.

External Merge Sort

- Divide-and-conquer sorting algorithm that splits the data set into separate **runs** and then sorts them individually.
- **Phase 1 – Sorting**
  - Sort blocks of data that fit in main-memory and then write back the sorted blocks to a file on disk.
- **Phase 2 – Merging**
  - Combine sorted sub-files into a single larger file.

2-Way External Merge Sort

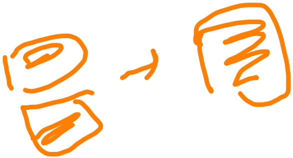
N pages = 100

- We will start with a simple example of a 2-way external merge sort.
  - "2" represents the number of runs that we are going to merge into a new run for each pass.
- Data set is broken up into **N** pages.
- The DBMS has a finite number of **B** buffer pages to hold input and output data.

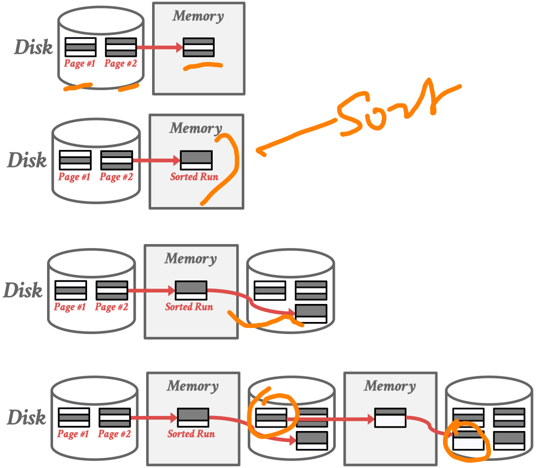
B buffer pages = 10

2-Way External Merge Sort

- **Pass 0**
  - Read every **B** pages of the table into memory
  - Sort pages into runs and write them back to disk.
- **Passes 1,2,3,. . .**
  - Recursively merge pairs of runs into runs **twice** as long.
  - Use three buffer pages (2 for input pages, 1 for output).



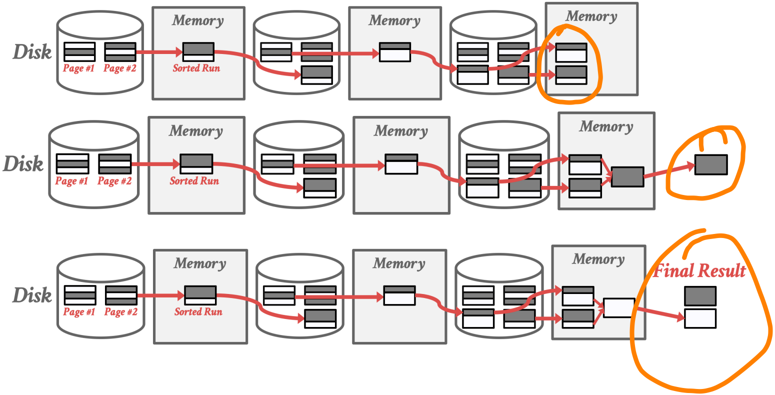
2-Way External Merge Sort



Sort

Sorting

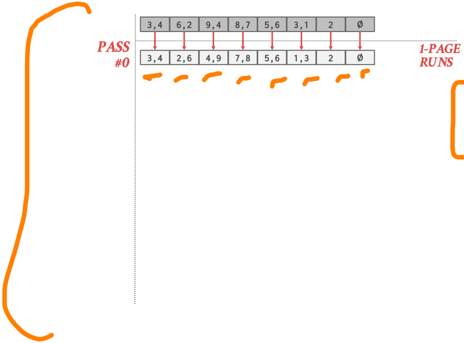
2-Way External Merge Sort

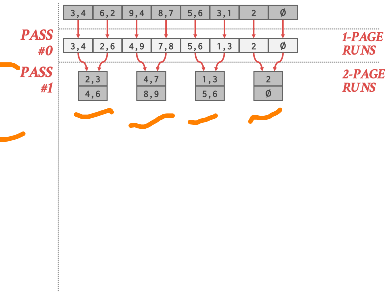


2-Way External Merge Sort

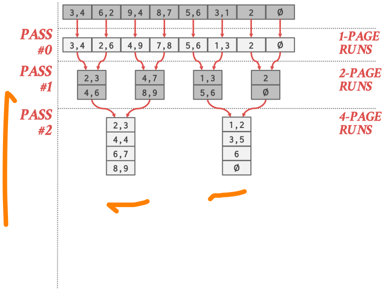
- In each pass, we read and write each page in file.
- Number of passes = $1 + \lceil \log_2 N \rceil$
- Total I/O cost = 2N x (Number of passes)

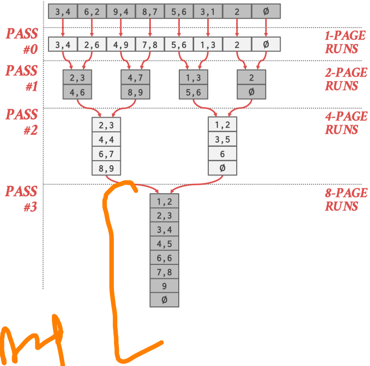
2-Way External Merge Sort





2-Way External Merge Sort





Sort
file

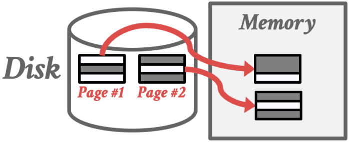
2-Way External Merge Sort

- This algorithm only requires three buffer pages to perform the sorting (**B=3**).
- But even if we have more buffer space available (**B>3**), it does not effectively utilize them.

Double Buffering Optimization

fetching ahead of time

- Prefetch the next run in the background and store it in a second buffer while the system is processing the current run.
  - Reduces the wait time for I/O requests at each step by continuously utilizing the disk.



CPU | I/O

General External Merge Sort

- **Pass 0**
  - Use **B** buffer pages.
  - Produce **N / B** sorted runs of size **B**
- **Pass 1,2,3,. . .**
  - Merge **B-1** runs (*i.e.*, K-way merge).
- Number of passes = $1 + \lceil \log_{B-1} N/B \rceil$
- Total I/O Cost = 2N x (Number of passes)

K-Way Merge Algorithm

- Input: **K** sorted sub-arrays
- Output: 1 sorted array
    - Efficiently compute the minimum element of all **K** sub-arrays.
    - Repeatedly transfer that element to output array
- Internally maintain a heap to efficiently compute minimum element.

Example

Sorting

N B

- Sort 108 pages with 5 buffer pages: **N=108**, **B=5**
  - Pass 0: **N / B** = 108 / 5 = 22 sorted runs of 5 pages each (last run is only 3 pages).
  - Pass 1: **N′ / B-1** = 22 / 4 = 6 sorted runs of 20 pages each (last run is only 8 pages).
  - Pass 2: **N″ / B-1** = 6 / 4 = 2 sorted runs, first one has 80 pages and second one has 28 pages.
  - Pass 3: Sorted file of 108 pages.
- $1+\log_{B-1} N/B = 1+\lceil \log_4 22 \rceil = 1 + \lceil 2.229 \rceil = 4$ passes

Merging.

- Table Heap
- B+Tree

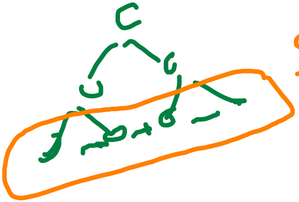
Tree-based Sorting

Using B+Trees for Sorting

Emp(C, ID)

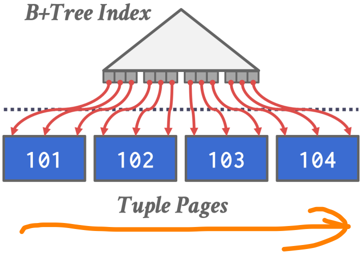
- If the table that must be sorted already has a B+Tree index on the sort attribute(s), then we can use that to accelerate sorting.
- Retrieve tuples in desired **sort order** by simply traversing the **leaf pages** of the tree.
- Cases to consider:
  - Clustered B+Tree
  - Unclustered B+Tree

Range Scan



Case 1 – Clustered B+Tree

- Traverse to the left-most leaf page, and then retrieve tuples from all leaf pages.
- This is always better than external sorting because there is no computational cost and all disk access is sequential.

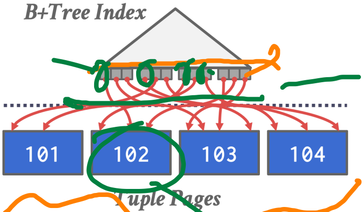


Case 2 – Unclustered B+Tree

Secondary Index

EMP_CITY

- Chase each pointer to the page that contains the data.
- This is almost always a bad idea. In general, one I/O per data record.



Sorted

once

EMP-ID

10

20

30

40

Value

3,

5,

10,

5,

5

Stable Key: EMP-ID

Aggregation

Aggregation

GROUP BY

- Collapse multiple tuples into a single scalar value.
- Two implementation choices:
  - ✓ Sorting
  - ✓ Hashing

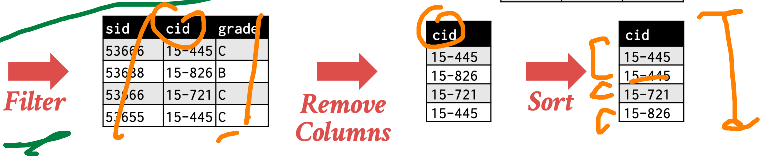
Sorting Aggregation

Sorting

```
SELECT DISTINCT cid
  FROM enrolled
 WHERE grade IN ('B','C')
 ORDER BY cid
```

enrolled(sid,cid,grade)

sid	cid	grade
53666	15-445	C
53688	15-721	A
53688	15-826	B
53666	15-721	C
53655	15-445	C



Filter

sid	cid	grade
53666	15-445	C
53688	15-826	B
53666	15-721	C
53655	15-445	C

Remove Columns

cid
15-445
15-826
15-721
15-445

Sort

cid
15-445
~~15-445~~
15-721
15-826

σ

Row Store

Sorting Aggregation

Grouping — HAVING n=0

Sorting — WHERE n=1

```
SELECT DISTINCT cid
  FROM enrolled
 WHERE grade IN ('B','C')
 ORDER BY cid
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enrolled(sid,cid,grade)

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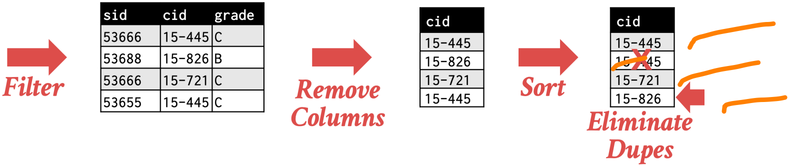
53666	15-721	C
53655	15-445	C



Filter

sid	cid	grade
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53688	15-826	B
53666	15-721	C
53655	15-445	C

Remove Columns

cid
15-445
15-826
15-721
15-445

Sort

cid
15-445
~~15-445~~
15-721
15-826

Eliminate Dupes

Alternatives to Sorting

n tuples

- What if we **do not** need the data to be ordered?
  - Forming groups in `GROUP BY` (no ordering)
  - Removing duplicates in `DISTINCT` (no ordering)
- Hashing is a better alternative in this scenario.
  - Only need to remove duplicates, no need for ordering.
  - May be computationally cheaper than sorting.

S: $O(n \lg n)$

H: $O(n)$

$O(1) \times N$

Hashing Aggregate

temporary

- Populate an **ephemeral hash table** as the DBMS scans the table.
- For each record, check whether there is already an entry in the hash table:
  - `GROUP BY`: Perform aggregate computation.
  - `DISTINCT`: Discard duplicates.
- If everything fits in memory, then it is easy. in-memory
- If the DBMS must spill data to disk, then we need to be smarter.

larger-than-memory

External Hashing Aggregate

hash (emp-city)

- **Phase 1 – Partition**
  - Divide tuples into buckets based on hash key.
  - Write them out to disk when they get full.
- **Phase 2 – ReHash**
  - Build in-memory hash table for each partition and compute the aggregation.

Phase 1 – Partition

- Use a hash function h_1 to split tuples into partitions on disk.
  - We know that all matches live in the same partition.
  - Partitions are **spilled** to disk via output buffers.
- Assume that we have **B** buffers.
- We will use **B-1** buffers for the partitions and **1** buffer for the input data.

B - 1 partitions.

Phase 1 – Partition

```
SELECT DISTINCT cid
  FROM enrolled
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```

enrolled(sid,cid,grade)

sid	cid	grade
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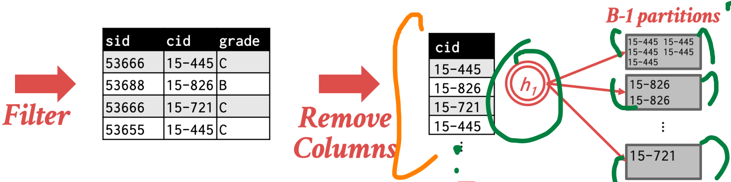
53666	15-721	C
53655	15-445	C



Filter

sid	cid	grade
53666	15-445	C
53688	15-826	B
53666	15-721	C
53655	15-445	C

Remove Columns

cid
15-445
15-826
15-721
15-445

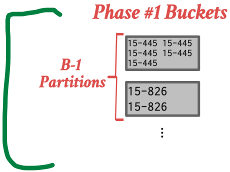
Phase 2 – ReHash

- For each partition on disk:
    - Read it into memory and build an in-memory hash table based on a second hash function h_2.
    - Then go through each bucket of this hash table to bring together matching tuples.
- This assumes that each partition fits in memory.

[illegible]
Chicago -

Phase 2 – ReHash

```
SELECT DISTINCT cid
  FROM enrolled
 WHERE grade IN ('B','C')
```



Phase #1 Buckets

B-1 Partitions

15-445 15-445
15-445 15-445
15-445

15-826
15-826

⋮

enrolled(sid,cid,grade)

sid	cid	grade
53666	15-445	C
53688	15-721	A
53688	15-826	B
53666	15-721	C
53655	15-445	C

Phase 2 – ReHash

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SELECT DISTINCT cid
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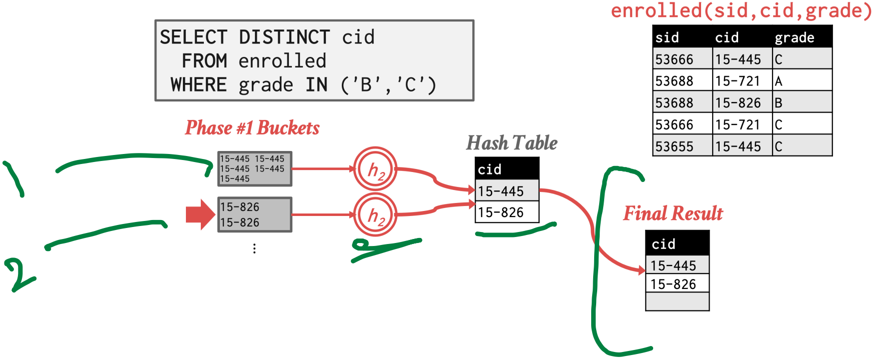

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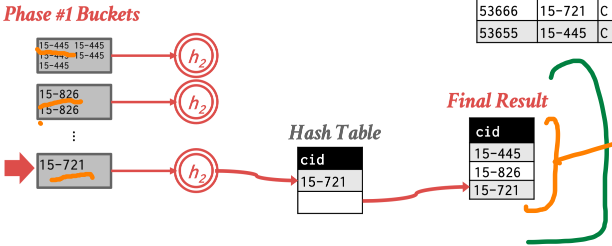
53688	15-721	A
53688	15-826	B
53666	15-721	C
53655	15-445	C



3

not Sorted

Hashing Summarization

City Name

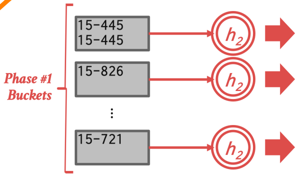
Total Salary, # of emps.

- During the `ReHash` phase, store pairs of the form (`GroupKey` ⟶ `RunningVal`)
- When we want to insert a new tuple into the hash table:
  - If we find a matching `GroupKey`, just update the `RunningVal` appropriately
  - Else insert a new `GroupKey` ⟶ `RunningVal`

Hashing Summarization

```
SELECT cid, AVG(s.gpa)
  FROM student AS s, enrolled AS e
 WHERE s.sid = e.sid
 GROUP BY cid
```

ORDER BY



Running Totals

AVG(col) → (COUNT,SUM)
MIN(col) → (MIN)
MAX(col) → (MAX)
SUM(col) → (SUM)
COUNT(col) → (COUNT)

Hash Table

key	value
15-445	(2, 7.32)
15-826	(1, 3.33)
15-721	(1, 2.89)

# of student

GPA

Final Result

cid	AVG(gpa)
15-445	3.66
15-826	3.33
15-721	2.89

S + A

Conclusion

Conclusion

requirements of
parent operator

- Choice of sorting vs. hashing is subtle and depends on optimizations done in each case.
- Next Class
  - Nested Loop Join
  - Sort-Merge Join
  - Hash Join