

Convex body  $K \subseteq \mathbb{R}^n$  in  $n$ -dimensions.

Given  $K$  by an oracle (membership):

For  $x \in \mathbb{R}^n$ , the oracle states whether  $x \in K$  or  $x \notin K$ .

Let  $B_n = B_n(0, 1)$  denote the unit ball at the origin in  $n$ -dimensions.

Assume  $B_n \subseteq K$  &  $K \subseteq DB_n$  for a given  $D > 1$ ,  
So  $D$  is the "diameter."

Measure running time by # of oracle calls &  
assume infinite precision since working with real numbers.

Sampling Problem: Generate a point  $x$  U.a.r. from  $K$ .

Counting problem: Compute  $\text{vol}(K)$ .

This  $\uparrow$  is #P-hard [Dyer-Frieze '88],  
hence can't hope to compute exactly in poly-time  
So aim for an FPRAS.

First FPRAS presented by [Dyer, Frieze, Kannan '91].

Latest is [Lovász, Vempala '03] which gives an  $O^*(n^4)$  time  
(i.e., oracle calls) for volume estimation.

Volume via sampling:

Given a sampling algorithm, how do we estimate  $\text{vol}(K)$ ?

Know  $B_n \subseteq K \subseteq DB_n$ .

Define sequence  $K_0 \subseteq K_1 \subseteq \dots \subseteq K_m$  for  $m = \Theta(n \log D)$ ,

where  $K_0 = B_n$  &  $K_m = K$ .

Also want that:  $\text{vol}(K_i) \leq 2 \text{vol}(K_{i-1})$ , i.e.,  $\frac{1}{2} \leq \frac{\text{vol}(K_{i-1})}{\text{vol}(K_i)} \leq 1$ .

Hence, if we generate  $x \in_R K_i$  then  $\Pr(x \in K_{i-1}) = \frac{\text{vol}(K_{i-1})}{\text{vol}(K_i)} \geq \frac{1}{2}$ .

Simplest approach:  $K_i = K \cap (2^{i/n} B_n)$ .

Claim:  $\text{vol}(K_i) \leq 2 \text{vol}(K_{i-1})$ .

Proof:  $K_i = 2^{i/n} B_n \cap K \subseteq 2^{i/n} (2^{(i-1)/n} B_n \cap K) = 2^{i/n} K_{i-1}$ .

By Chebyshev's inequality, need  $O(m/\epsilon^2)$  samples per  $K_i$ .

&  $O(m^2/\epsilon^2)$  in total to estimate  $(1 \pm \epsilon) \text{vol}(K)$

with prob.  $\geq 3/4$  using:

$$\text{vol}(K) = \frac{\text{vol}(K_m)}{\text{vol}(K_{m-1})} \times \frac{\text{vol}(K_{m-1})}{\text{vol}(K_{m-2})} \times \dots \times \frac{\text{vol}(K_1)}{\text{vol}(K_0)} \times \text{vol}(B_n)$$

$$\& \text{vol}(B_{2n}) = \pi^n / n!$$

Better approach in [Lovasz-Vempala '03]: Use exponential distributions.

How to sample?

Ball walk: Two versions - lazy & speedy.

$$\text{Let } \delta = \Theta\left(\frac{1}{\sqrt{n}}\right)$$

Lazy walk: (this is the algorithm used)

From  $X_t \in K$ :

1. With prob.  $\frac{1}{2}$ , set  $X_{t+1} = X_t$ .

2. Else:

- Choose  $X'$  uniformly at random from  $B(X_t, \delta)$

- If  $X' \in K$ , set  $X_{t+1} = X'$   
else  $X_{t+1} = X_t$ .

Speedy walk: (just used in the proof)

From  $X_t \in K$ :

1. Choose  $X_{t+1}$  uniformly at random from  $B(X_t, \delta) \cap K$ .

How do you implement speedy walk?

For  $x \in K$ , let  $l(x) = \frac{\text{vol}(B(x, \delta) \cap K)}{\text{vol}(B(x, \delta))}$

To implement speedy: do lazy until get a point  $x \in K$ ; this takes  $1/l(x)$  steps in expectation.

What's the invariant measure (i.e., stationary distribution)?

Statespace is uncountably infinite so standard results don't apply.

For lazy walk, invariant measure  $\mu$  is uniform:

for  $A \subseteq K$ ,  $\mu(A) = \frac{\text{vol}(A)}{\text{vol}(K)}$

(this is unique, see Thm. 2.1 in [Vempala '05] MSR notes or Section 1 in [Lovász-Simonovits '93]

For speedy walk,

$\mu(A) = \frac{\int_A l(x) dx}{L}$  where  $L = \int_K l(x) dx$

(5)

For speedy walk,  $\Phi \geq \frac{c\delta^2}{D^2 n}$  for constant  $c > 0$ .

This is the "right" answer b/c think of an unbiased walk on the integers  $0, 1, \dots, D$ , which takes  $\mathcal{O}(D^2)$  time to mix. And here the walk moves  $\approx \frac{\delta}{\sqrt{n}}$  in 1 direction, and hence rescale to  $\frac{D\sqrt{n}}{\delta}$  & this becomes  $\mathcal{O}\left(\frac{D^2}{\delta}\right)$ .

This implies rapid mixing for lazy walk from a "warm start" which is an initial distribution  $w$  where  $\forall A \subseteq K$ :  $w(A) \leq 2\mu(A)$ , which we obtain by using a random sample from  $K_{i-1}$  as the initial distribution on  $K_i$ .

— for a warm start unlikely to start in a "corner" & unlikely to ever get into a corner.

(6)

To bound conductance look at the following generalization:

$$\text{let } \lambda = \min_{\substack{f: K \rightarrow \mathbb{R}: \\ f \text{ is not constant}}} \frac{E_{\mu}(f, f)}{\text{Var}_{\mu} f}$$

$$\text{when } f = 1(S) = \begin{cases} 1 & \text{if } x \in S \\ 0 & \text{if } x \notin S \end{cases} \quad \text{for } S \subset K$$

$$\text{then } \frac{E_{\mu}(f, f)}{\text{Var}_{\mu} f} = \Phi(S).$$

for discrete-time MC on  $(\Omega, \mathcal{P}, \pi)$ :

$$\text{Var}_{\pi} f = \sum_{x \in \Omega} \pi(x) f(x)^2 - \left( \sum_{x \in \Omega} \pi(x) f(x) \right)^2$$

$$= \sum_{x \in \Omega} \pi(x) f(x)^2 - \left( \sum_{x \in \Omega} \pi(x) f(x) \right)^2$$

$$= \sum_{x \in \Omega} \pi(x) f(x)^2 \sum_{y \in \Omega} \pi(y) - \sum_x \pi(x) f(x) \sum_y \pi(y) f(y)$$

$$= \sum_{x, y \in \Omega} \pi(x) \pi(y) (f(x)^2 - f(x) f(y))$$

$$= \frac{1}{2} \sum_{x, y \in \Omega} \pi(x) \pi(y) (f(x) - f(y))^2$$

(7)

Hence,  $\text{Var}_\pi f$  measures how  $f$  varies globally: across all Pairs  $x, y$ .

instead  $\sum_\pi (f, f)$  measures the "local" variation of  $f$  with respect to transitions  $P$ :

$$\text{Thus, } \sum_\pi (f, f) = \frac{1}{2} \sum_{x, y \in \Omega} \pi(x) P(x, y) (f(x) - f(y))^2$$

In fact,  $\lambda = (\text{gap})^{-1}$  where  $\text{gap} = 1 - \max\{\lambda_i, P_{ii}\}$   
for eigenvalues  $\lambda_0 = 1 > \lambda_1 \geq \dots \geq \lambda_n > -1$   
of  $P$

When  $f = 1(s)$  then:

$$\sum_\pi (f, f) = \frac{1}{2} \sum_{\substack{x \in S, y \notin S \\ \text{or} \\ x \notin S, y \in S}} \pi(x) P(x, y) = \sum_{x \in S, y \notin S} \pi(x) P(x, y)$$

for reversible MC

$$\Delta \text{Var}_\pi f = \frac{1}{2} \sum_{\substack{x \in S, y \notin S \\ \text{or} \\ x \notin S, y \in S}} \pi(x) \pi(y) = \pi(S) \pi(\bar{S}).$$

Let  $\lambda = \min_{f: K \rightarrow \mathbb{R}} \frac{E_u(f, f)}{\text{Var}_u f}$

Note: can restrict to  $f$  s.t.  $\bar{f} = E_u f = 0$

by shifting (doesn't change  $E_u(f)$  or  $\text{Var}_u f$ )

For  $A \subseteq K$ ,

$$P(x, A) = \Pr(X_1 \in A \mid X_0 = x)$$

for speedy walk,

$$P(x, A) = \frac{\text{vol}(B(x, \delta) \cap A)}{\text{vol}(B(x, \delta) \cap K)}$$

for  $y \in B(x, \delta) \cap K$ ,

$$P(x, dy) = \frac{dy}{\text{vol}(B(x, \delta) \cap K)}$$



For  $x \in K$ , let

$$\begin{aligned}
 h(x) &= \frac{1}{2} \int_K F(x, dy) (f(x) - f(y))^2 \\
 &= \frac{1}{2 \text{vol}(B(x, \delta) \cap K)} \int_{B(x, \delta) \cap K} (f(x) - f(y))^2 dy
 \end{aligned}$$

Assume WOLOG that  $\bar{f} = E_{\mu} f = 0$

& then  $\text{Var}_{\mu} f = E_{\mu} f^2 = \int_K f^2 d\mu$

&  $E_{\mu}(f, f) = \int_K h d\mu$

We want to show: for all  $f$  s.t.  $\bar{f} = 0$ ,

$$\frac{\int_K h d\mu}{\int_K f^2 d\mu} \geq \lambda \text{ for } \lambda = \frac{c\delta^2}{D^n}$$

This is the argument from

[KLS] = [Kannan, Lovász, Simonovits '95]

Reduction to "needle-like" case:

(10)

Suppose  $\exists f$  where  $\frac{\int_K h du}{\int_K f^2 du} < \lambda$  &  $\int_K f du = 0$

We'll show that there is a convex  $K_1 \subseteq K$  where:

$$\frac{\int_{K_1} h du}{\int_{K_1} f^2 du} < \lambda \text{ \& \ } \int_{K_1} f du = 0$$

and  $K_1 \subseteq [0, D] \times [0, \epsilon]^{n-1}$

for any  $\epsilon > 0$  (can make arbitrarily small)

(Roughly, if there's a violating  $f$   
then there's a violating  $f$  for a  
1-dimensional problem)

(11)

# Key geometric facts:

A: For convex  $P \subset \mathbb{R}^2$ ,

$\exists$  point  $x \in P$  s.t. every line  $l$  through  $x$ ,

the two sides  $P^+ = P \cap l^+$  &  $P^- = P \cap l^-$

have  $\text{area}(P^+) \geq \frac{\text{area}(P)}{3}$  &  $\text{area}(P^-) \geq \frac{\text{area}(P)}{3}$

B: For convex  $P \subset \mathbb{R}^2$  of area  $A$ ,

$$\text{width}(P) \leq \sqrt{2A}$$

where  $\text{width} = \min \text{distance}(l, l')$

$l, l' \neq$

$l \& l'$  are parallel

and sandwich  $P$ .

(12)

Suppose for  $j \geq 2$  we have convex  $K_j \subseteq \mathbb{R}^n$

where  $K_j \subseteq [0, D]^j \times [0, \epsilon]^{n-j}$  &  $\int_{K_j} f du = 0$

$$\& \frac{\int_{K_j} h du}{\int_{K_j} f^2 du} < \lambda$$

Then we can reduce  $j$  by 1 as follows:

Base case:  $j = n$ :

For  $2 \leq j \leq n$ : these  $j$  are "fat" dimensions.

Take 1<sup>st</sup> 2 fat coordinates.

Look at projection of  $K_j$  onto them, call it  $P$ .

Take  $x \in P$  from  $A$ .

Take  $(n-1)$ -dimensional plane  $G$  through  $x$  whose normal lies in  $P$ .

We know  $\int_{K_j} f du = 0$

So  $\int_{K_j \cap G^+} f du + \int_{K_j \cap G^-} f du = 0$

either these  $\uparrow$   $\nearrow$  are = or one is  $> 0$  & other is  $< 0$ .

If flip  $G$ , then signs  $\nearrow$  flip.

Now rotate  $G$  & it changes continuously so at some point it changes from + to -

So they are = for some  $H$ .

Thus,  $\int_{K_j \cap H^+} f du = \int_{K_j \cap H^-} f du = 0$

We know  $\int_{K_j} h du < \lambda \int_{K_j} f^2 du$

So either:  $\int_{K_j \cap H^+} h du < \lambda \int_{K_j \cap H^+} f^2 du$

or  $\int_{K_j \cap H^-} h du < \lambda \int_{K_j \cap H^-} f^2 du$ .

Take the violating set  $\mathcal{Q}$  repeat

Each time  $\text{area}(P)$  decreases by  $\geq 2/3$

So eventually  $\text{area}(P) \leq \frac{1}{2}\epsilon^2$

then by B, width of  $S$  is  $\leq \epsilon$ .

Finally, rotate  $\mathcal{Q}$  one of these 2 dimensions is now of width  $\leq \epsilon$ .