

Glauber Dynamics for the Ising Model

①

Ising Model

$$G = (V, E)$$



$$\Omega = \{+1, -1\}^{|V|}$$

For $\sigma \in \Omega$,

$$\mu(\sigma) = \frac{1}{Z} \exp\left(\beta \cdot \sum_{u,v} \mathbb{1}(\sigma_u = \sigma_v)\right)$$

$$Z = \sum_{\sigma \in \Omega} \exp\left(\beta \sum_{u,v} \mathbb{1}(\sigma_u = \sigma_v)\right)$$

Glauber dynamics

$$X_t \in \Omega$$

1) Pick $v \in V$ u.a.r.

2) Update the spin of v with a sample from:

$$\mu(\cdot \mid X_t(v, w))$$

Let $N_v^+(X_t) = \#$ of "+" in neighborhood of v in X_t .

$N_v^-(X_t) = \#$ of "-" " " "

Then,

$$\mu(v = + \mid X_t(v, w)) = \frac{e^{\beta N_v^+(X_t)}}{e^{\beta N_v^+(X_t)} + e^{\beta N_v^-(X_t)}}$$

$$\mu(v=- | X_{\mathbb{Z}}(v,w)) = \frac{e^{\beta N_v^-(x_{\mathbb{Z}})}}{e^{\beta N_v^-(x_{\mathbb{Z}})} + e^{\beta N_v^+(x_{\mathbb{Z}})}}$$

Proof:

$$\mu(v=+ | X_{\mathbb{Z}}(v,w)) = \frac{\mu(v=+, X_{\mathbb{Z}}(v,w))}{\mu(v=+, X_{\mathbb{Z}}(v,w)) + \mu(v=-, X_{\mathbb{Z}}(v,w))}$$

⇒ $\mathbb{S}D$ dynamics is ergodic and reversible w.r.t. μ .

$$\sigma_v(u) = \sigma(u) \quad \forall u \neq v \quad \text{and} \quad \sigma_v(v) = -\sigma(v)$$

→ Typo: sigma_v should be sigma^v

$$\mu(\sigma) \cdot P(\sigma, \sigma^{(v)}) = \mu(\sigma^{(v)}) \cdot P(\sigma^{(v)}, \sigma)$$

$$= \frac{1}{Z} \exp \left[\beta \sum_{\substack{u,v \\ u,v \neq v}} \mathbb{1}(\sigma_u = \sigma_w) \right] \exp \left[\beta \cdot N_v^{\sigma(v)}(\sigma) \right] \cdot \frac{e^{\beta N_v^{\sigma(v)}(\sigma)}}{e^{\beta N_v^{\sigma(v)}(\sigma)} + e^{\beta N_v^{-\sigma(v)}(\sigma)}}$$

$$= \frac{1}{Z} \exp \left[\beta \sum_{\substack{u,v \\ u,v \neq v}} \mathbb{1}(\sigma_u = \sigma_w) \right] \exp \left[\beta \cdot N_v^{-\sigma(v)}(\sigma) \right] \cdot \frac{\exp[\beta N_v^{\sigma(v)}(\sigma)]}{e^{\beta N_v^{\sigma(v)}(\sigma)} + e^{\beta N_v^{-\sigma(v)}(\sigma)}}$$

$$= \mu(\sigma^{(v)}) \cdot P_i[\sigma^{(v)}, \sigma]$$

Mixing TIME: $\beta_c = \ln(1 + \sqrt{2})$

(i) If $\beta < \beta_c$, then $T_{mix} = O(n \log n)$

(ii) If $\beta > \beta_c$, then $T_{mix} = e^{\Omega(\sqrt{n})}$

Intuition:

(i) Coupon collecting



Bottleneck btw. mostly "+" and mostly "-" configurations.

⇒ We will prove (i) today. (3)

⇒ In fact we'll prove something slightly weaker.

Thm 1: If $\beta < \beta_c$, then $T_{\text{mix}} = O(n \log^3 n)$.

[Proof due to Dyer, Sinclair, Vigoda, Weitz '04].

Lemma 2: Suppose $T_{\text{mix}}(\sqrt{n} \times \sqrt{n}) \leq g(n)$
where $g: \mathbb{N} \rightarrow \mathbb{R}^+$ is an increasing function.

Then,
 $T_{\text{mix}}(\sqrt{n} \times \sqrt{n}) \leq n \cdot g((\log n)^2)$.

Proof of Thm 1:

Suppose we know $T_{\text{mix}}(\sqrt{n} \times \sqrt{n}) \leq e^{cn}$.
Then $T_{\text{mix}}(\sqrt{n} \times \sqrt{n}) \leq n \cdot e^{c(\log n)^2} \leq e^{c_1(\log n)^2}$
 $T_{\text{mix}}(\sqrt{n} \times \sqrt{n}) \leq n \cdot e^{c(2 \log \log n)^2} \leq e^{c_2 \log n}$
 $\leq n^{c_2}$

$$T_{\text{mix}}(\sqrt{n} \times \sqrt{n}) \leq n \cdot (\log n)^{2c_2}$$

$$T_{\text{mix}}(\sqrt{n} \times \sqrt{n}) \leq n \cdot (\log n)^2 (\log \log n)^{c_3}$$

$$T_{\text{mix}}(\sqrt{n} \times \sqrt{n}) \leq n \cdot (\log n)^2 (\log \log n)^{c_4}$$

Proof of Lemma 2

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Idea: Design coupling s.t. after $T = n \cdot g(\log n^2)$

Steps:

$$\Pr[X_T \neq Y_T] \leq 1/4$$

and the result follows.

Identity Coupling:

⇒ Pick same ^(random) vertex v in both copies.

⇒ Use same random coin to sample from $\mathcal{U}(\cdot | X_t(v,v))$ and $\mathcal{U}(\cdot | Y_t(v,v))$.

• Suppose "+" \geq "-" (total order of the spins).

• Then $X_t \geq Y_t$ if $X_t(v) \geq Y_t(v) \forall v \in V$.
(partial order of the state space).

Key Fact: The identity coupling is monotone

That is, if $X_t \geq Y_t$, then

$$X_{t+1} \geq Y_{t+1}.$$

Proof: • Suppose v is the vertex chosen to be updated.

• Then X_t has at least the same number of "+" in the neighborhood of v as Y_t .

$$(N_v^+(X_t) \geq N_v^+(Y_t)).$$

Then

$$M(v=+ | X_t(v|v)) = \frac{e^{\beta N_t^+(x_t)}}{e^{\beta N_t^+(x_t)} + e^{\beta N_t^-(x_t)}}$$

$$= \frac{1}{1 + e^{\beta(N_t^-(x_t) - N_t^+(x_t))}} = \frac{1}{1 + e^{\beta[d - 2N_t^+(x_t)]}}$$

$$\geq \frac{1}{1 + e^{\beta[d - 2N_t^+(y_t)]}} = M(v=+ | Y_t(v|v)).$$

So, we need to prove (assuming $X_0 = \text{all "+"}$, $Y_0 = \text{all "-"}$)

$$\Pr[X_T \neq Y_T] \leq \frac{1}{4} \quad T = n \cdot g((\log n)^2).$$

$$\Pr[X_T \neq Y_T] \leq \sum_{v \in V} \Pr[X_T(v) \neq Y_T(v)] \quad (\text{union bound}).$$

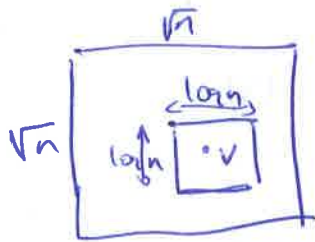
So, enough to show:

$$\Pr[X_T(v) \neq Y_T(v)] \leq \frac{1}{4n}.$$

(assuming $X_0 = \text{"+"}$ and $Y_0 = \text{"-"}$)

Let $\{Z_t^+\}$ be an auxiliary copy of the Markov^⑥ chain, such that:

- $Z_0^+ =$ all "+".
- $\{Z_t^+\}$ does not perform any updates outside of a box ~~of size $\log n$~~ centered at v of side length $\log n$.



Similarly define $\{Z_t^-\}$. ($Z_0^- =$ all minus and $\{Z_t^-\}$ does not update in box.)

Then, by monotonicity:

$$\Pr[X_t(v) \neq Y_t(v)] \leq \Pr[Z_t^+(v) \neq Z_t^-(v)]$$

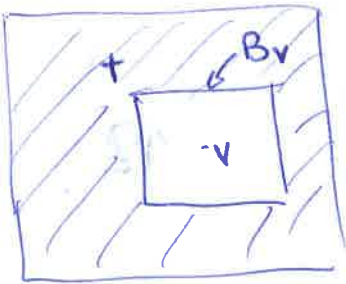
$$\text{Since } Z_t^- \subseteq Y_t \subseteq X_t \subseteq Z_t^+ \quad \forall t \geq 0$$

$$\begin{aligned} \text{Now, } \Pr[Z_t^+(v) \neq Z_t^-(v)] &= \Pr[Z_t^+(v) = +, Z_t^-(v) = -] \\ &= \Pr[Z_t^-(v) = - \mid Z_t^+(v) = +] \cdot \Pr[Z_t^+(v) = +] \\ &= (1 - \Pr[Z_t^-(v) = + \mid Z_t^+(v) = +]) \Pr[Z_t^+(v) = +] \\ &= \Pr[Z_t^+(v) = +] - \Pr[Z_t^-(v) = +] \end{aligned}$$

Then,

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$$\Pr[X_t(v) \neq Y_t(v)] \leq \Pr[Z_t^+(v) = +] - \Pr[Z_t^-(v) = +]$$



~~Proof~~

Idea: after $T = n g((\log n)^2)$ steps, how many hit B ?

$$\approx \frac{(\log n)^2}{n} \cdot n g((\log n)^2)$$

$$= g((\log n)^2) (\log n)^2,$$

which by assumption is much larger than

$$T_{\text{mix}}(\log n \times \log n) = g((\log n)^2).$$

\Rightarrow So, the GD is very well-mixed in B_v .

\Rightarrow The stationary measure in B_v is $\mu(\cdot | B_v^c = "+")$ for $\{Z_t^+\}$ and $\mu(\cdot | B_v^c = "-")$ for $\{Z_t^-\}$.

\Rightarrow Then,

$$\Pr[Z_T^+(v) = 0] \approx \mu(v = + | B_v^c = "+").$$

$$\Pr[Z_T^-(v) = 0] \approx \mu(v = + | B_v^c = "-").$$

\Rightarrow All that we need is then

$$\mu(v = + | B_v^c = "+") \approx \mu(v = - | B_v^c = "-").$$

⇒ This property holds for $\beta < \beta_c$:

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$$(*) \left| \mu(V=+ | B_v^c = "+") - \mu(V=+ | B_v^c = "-") \right|$$

$$\leq \exp(-\text{dist}(v, \partial B_v)) \approx \frac{1}{12n}$$

↳ [This property is called Strong Spatial Mixing]

⇒ Consequently, $\Pr[X_T(v) \neq Y_T(v)]$ should be "small". Let's quantify it:

$$\Pr[X_T(v) \neq Y_T(v)] \leq \Pr[Z_T^+(v) \neq "+"] - \Pr[Z_T^-(v) = "-"]$$

$$\leq \left| \Pr[Z_T^+(v) = "+"] - \mu(V=+ | B_v^c = "+") \right| \quad (A)$$

$$+ \left| \Pr[Z_T^-(v) = "+"] - \mu(V=+ | B_v^c = "-") \right| \quad (B)$$

$$+ \left| \mu(V=+ | B_v^c = "+") - \mu(V=+ | B_v^c = "-") \right| \quad (C)$$

⇒ By the Strong Spatial Mixing property (C) $\leq \frac{1}{12n}$.

⇒ Note that (A) $\leq \frac{1}{12n}$ also, since after

$$T = n \cdot g((\log n)^2) \quad \left\| Z_T^+ - \mu(\cdot | B_v^c = "+") \right\|_{TV} \leq \frac{1}{4(\log n)^2} \leq \frac{1}{12n}$$

Recall that if:

$$\|P^\tau(x_0, \cdot) - \mu\|_{TV} \leq \frac{1}{4},$$

then

$$\|P^{\tau k}(x_0, \cdot) - \mu\|_{TV} \leq \frac{1}{4^k}$$

Putting all these together.

$$\textcircled{A} + \textcircled{B} + \textcircled{C} \leq \frac{3}{12n} = \frac{1}{4n}$$

