

For an undirected $G=(V,E)$,

let $\mathcal{T}(G) = \text{spanning trees of } G$

$= \{S \subseteq E: S \text{ is acyclic \& connected}\}$

thus $|S| = |V| - 1$

Goal: sample uniformly from \mathcal{T} .

For undirected G , define \vec{G} to be the directed graph
where: for $(v,w) \in E$, add \vec{vw} & \vec{wv} to \vec{E}

For directed \vec{G} , for vertex $r \in V$,
arborescence rooted at r :

$S \subseteq E$ s.t. $|S| = |V| - 1$

& every $v \neq r$ has exactly 1 edge
directed away from v .

(i.e., it's a directed tree rooted at r ,
pointing ~~away from~~ towards r)

Let $\mathcal{T}_r(\vec{G}) = \text{set of arborescences rooted at } r$.

& $\mathcal{T}(\vec{G}) = \bigcup_{r \in V} \mathcal{T}_r(\vec{G}) = \text{all arborescences}$.

Observation: $|\mathcal{J}(G)| = |\mathcal{J}_r(\vec{G})|$

spanning trees of G = # arborescences of \vec{G} rooted at r .

Proof:

\Rightarrow for $S \in \mathcal{J}(G)$, form $S' \in \mathcal{J}_r(\vec{G})$ by pointing all edges ~~away from~~ towards r .

\Leftarrow for $S' \in \mathcal{J}_r(\vec{G})$, drop orientations & we have $S \in \mathcal{J}(G)$.

MC1 on $\mathcal{J}(\vec{G})$: (on all arborescences)

From $X_t \in \mathcal{J}(\vec{G})$ with root v .

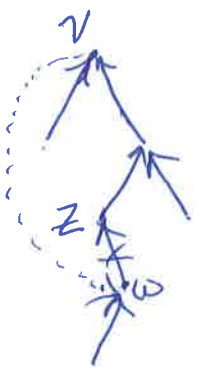
1. Choose a random ^(out) neighbor w of v (i.e., $\vec{vw} \in E$)

2. Let \vec{wz} be the unique edge away from w in X_t

3. Let $X' = X_t \cup \vec{vw} \setminus \vec{wz}$.

Note $X' \in \mathcal{J}(\vec{G})$ with root w

4. Let $X_{t+1} = \begin{cases} X' & \text{with prob. } 1/2 \\ X_t & \text{otherwise.} \end{cases}$



③

If G is \mathcal{Q} -regular then MC1 is symmetric & thus π is uniform over $\mathcal{Z}(G)$.

Fix root $r \in V$.

Here is a new MC call it MC_r on $\mathcal{Z}_r(G)$.

From $Y_t \in \mathcal{Z}_r(G)$:

1. Run MC1 with $X_0 = Y_t$

Stop at first time t' where

$$X_{t'} \in \mathcal{Z}_r(G)$$

(i.e., 1st time when root returns to r .)

2. Set $Y_{t+1} = X_{t'}$

Note: this MC_r is symmetric & has uniform over $\mathcal{Z}_r(G)$ stationary distribution.

What's the mixing time?

Consider ^{lazy} random walk on G .

From $v \in V$, with prob. $\frac{1}{2}$ stay at v ,

& with prob. $\frac{1}{2}$ move to a random neighbor.

Cover time = # of steps until random walk visits every vertex at least once from worst starting state.
// T_{cover}

Lemma: ~~$E[T_{mix}] \leq$~~

$$T_{mix} \leq 4E[T_{cover}]$$

& for all G ,

$$n \log n \leq E[T_{cover}] \leq O(nm).$$

Let's prove the mixing time.

Coupling: For a pair $Y_+, Y'_+ \in \mathcal{J}_r(\vec{G})$
 we'll define a coupling $(Y_+, Y'_+) \rightarrow (Y_{++}, Y'_{++})$
 where:

for all $Y_0, Y'_0 \in \mathcal{J}_r(\vec{G})$

$$\Pr(Y_T \neq Y'_T) \leq \frac{1}{4}$$

for $T = 4E[T_{\text{cover}}]$.

For Y_+, Y'_+ , let

$$\begin{aligned} H(Y_+, Y'_+) &= |Y_+ \setminus Y'_+| \\ &= \# \text{ edges in } Y_+ \text{ \& not in } Y'_+ \end{aligned}$$

thus if $H(Y_+, Y'_+) = 0$ then $Y_+ = Y'_+$.

What's coupling of $(Y_t, Y'_t) \rightarrow (Y_{t+1}, Y'_{t+1})$?

$Y_t \rightarrow Y_{t+1}$ uses "excursion"

$X_0 \rightarrow X_1 \rightarrow X_2 \rightarrow \dots \rightarrow X_t$

root: $r \rightarrow v_1 \rightarrow v_2 \rightarrow \dots \rightarrow r$
chosen edge: $(r, v_1), (v_1, v_2), \dots, (v_{t-1}, r)$

for $Y'_t \rightarrow Y'_{t+1}$ use same sequence of edges

Then they share these edges, namely

for v on this excursion,
let i be the last time.

Then, $\overrightarrow{v_i v_{i+1}} \in Y_{t+1} \& \# \overrightarrow{v_i v_{i+1}} \in Y'_{t+1}$

Once every vertex is visited at least once
on an excursion then $Y_t = Y'_t$.

The excursions are just a lazy random walk
on G .

The time to visit every vertex at least once is the cover time.

$$\text{Let } T = 4E[T_{\text{cover}}].$$

$$\begin{aligned} \text{Then, } \Pr(Y_T \neq Y'_T) &\leq \Pr(T_{\text{cover}} > T) \\ &= \Pr(T_{\text{cover}} > 4E[T_{\text{cover}}]) \\ &\leq \frac{1}{4}. \end{aligned}$$

□