

# Volume Computation

Thursday, November 9, 2017 11:11 AM

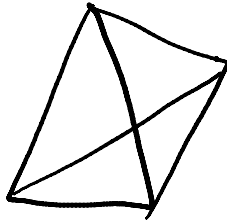
Yuler

Basic Problem: Find the volume of a given set.  
in  $\mathbb{R}^n$   
classical, ancient.

e.g. pyramids (Egypt), wine barrels (Greece)

Some shapes have nice formulas

e.g. Simplex = conv  $\{x_0, \dots, x_n\}$


$$\text{vol} = \frac{1}{n!} \left| \det \begin{pmatrix} 1 & 1 & \dots & 1 \\ x_0 & x_1 & \dots & x_n \end{pmatrix} \right|$$

$$\text{Ball}(r) = c_n \cdot r^n$$

$$\text{Ellipsoid} = \{x : x^T A^{-1} x \leq 1\}$$

$$\text{vol} = \sqrt{\det(A)} \cdot \text{vol}(B_n).$$

Parallelepiped

⋮

Not a long list.

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0 1 1... 1 2

1000 ...

For a polytope  $P = \{x : Ax \geq b\}$

$A, b$  rational. How to compute the volume?

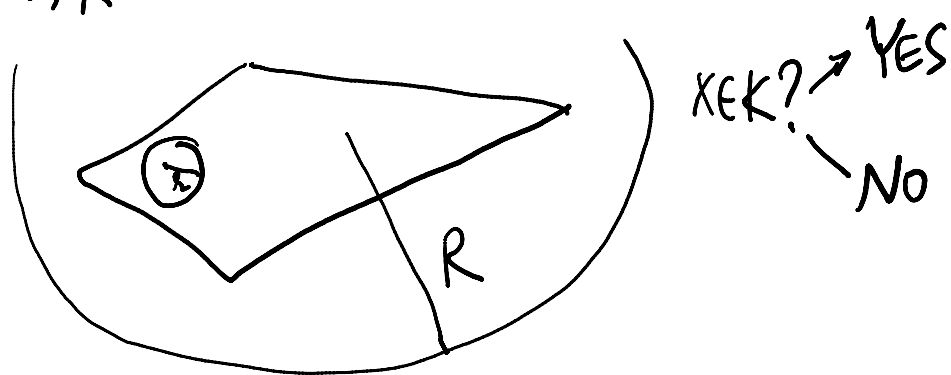
[Dyer-Frieze] Thm. #P-hard.

We can hope to approximate it.

But let's generalize first.

Input: body  $K$  given by a  
"well-guaranteed" membership oracle

$$\epsilon, R : x_0 + \epsilon B_n \subseteq K \subseteq R B_n$$



$$\epsilon > 0.$$

Output:  $V$  s.t.  $(1-\epsilon) \text{Vol}(K) \leq V \leq (1+\epsilon) \text{Vol}(K)$

Even more general:

Input:  $f: \mathbb{R}^n \rightarrow \mathbb{R}_+$  nonnegative function

$\int f < \infty$  integrable

given by a "well-guaranteed" function oracle

$x \rightarrow \boxed{\phantom{x}} \rightarrow f(x) \quad \epsilon > 0$

Output:  $F$  s.t.  $(1-\epsilon) \int f \leq F \leq (1+\epsilon) \int f$ .

Eg. What is the Gaussian measure of a polytope?

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Need to make assumptions even to be able to approximate.

$K$ -convex

$f$ -logconcave (we'll see later)

(common generalization of  $\chi_K$  and  $e^{-\frac{\|x\|^2}{2}}$ )

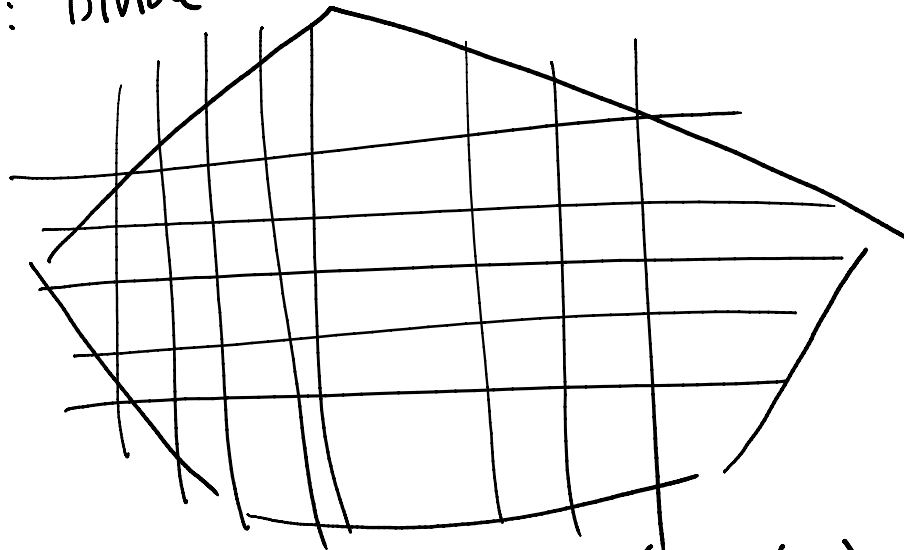
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Problem. Approximate volume of a given convex body

" " " " + oracle arithmetic complexity.

Complexity : # calls to oracle, arithmetic complexity.

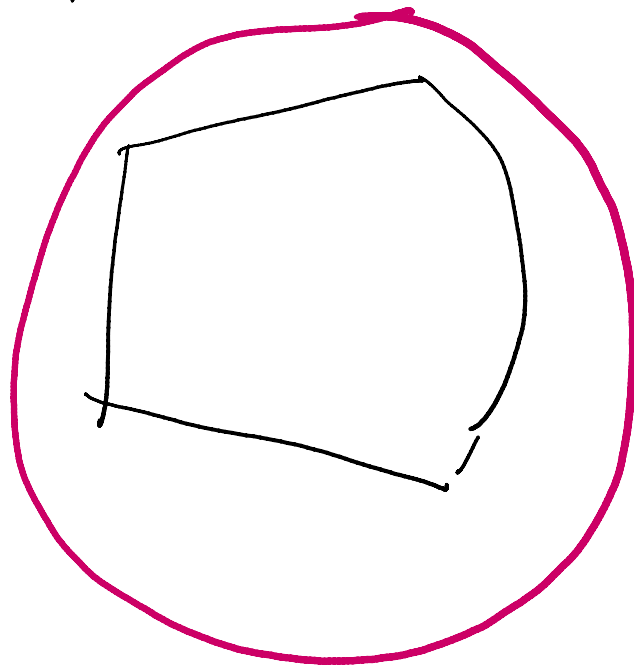
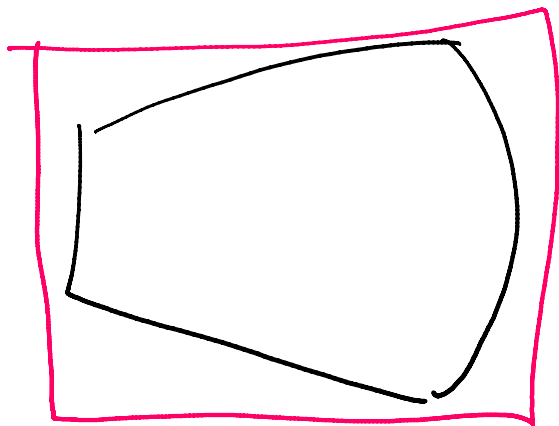
Attempt #1 : Divide-and-conquer



For small enough cubes,  $\text{vol}(\cup \text{cubes}) \approx \text{vol}(K)$ .

Difficulty?

Attempt #2 : Enclose in simple shape



Difficulty?



Difficulty?

Attempt #3. Divide polytope  $P$  into simplices.

Difficulty?

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Thm. [E, BF] For Any deterministic algorithm that uses at most  $n^a$  oracle calls to estimate  $A, B$  s.t.  $A \leq \text{vol}(K) \leq B$ ,  $\exists K$  s.t.

$$\frac{B}{A} > \left( \frac{cn}{a \log n} \right)^{\frac{n}{2}}$$

Thm. [BF]  $2^{cn}$  deterministic approximation is  $\exists c, c'$  s.t. not possible in  $2^{c'n}$  time.

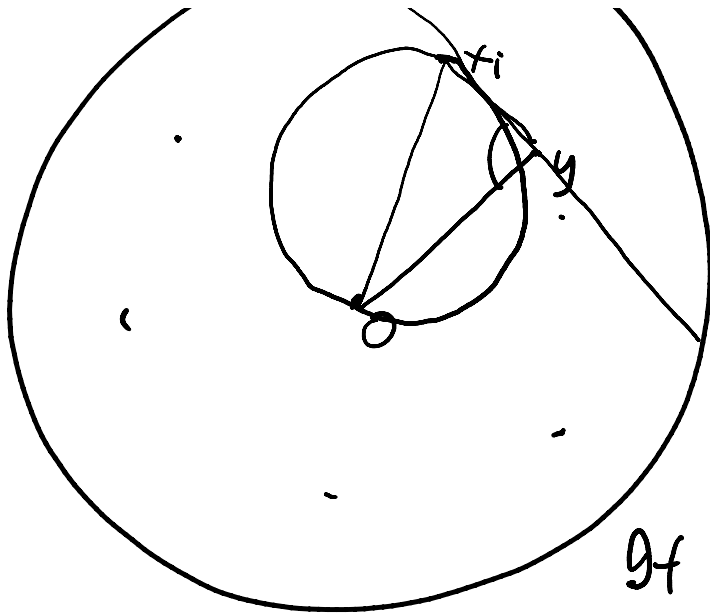
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Why difficult?



$x_1 \dots x_m$

Lemma  $\text{Conv} \{x_1 \dots x_m\} \subseteq \cup B_i$



$$1. \subseteq \cup B_i$$

$$2. \text{Vol}(\cup B_i) \leq \frac{m}{2^n}.$$

PF. Take any  $y \in \text{conv}\{x_1, \dots, x_m\}$

If  $y \notin B_i$ ,

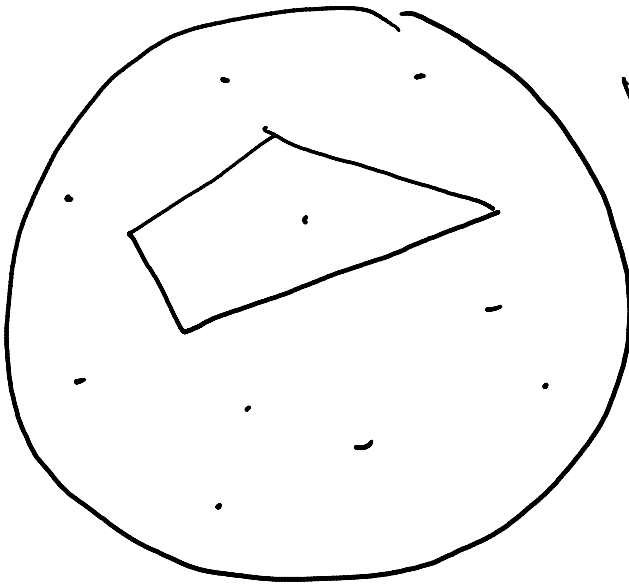
then angle at  $y > \frac{\pi}{2}$

$$\Rightarrow y^T x_i < y^T y$$

But this holds for all  $x_i$  !

So  $y$  can be separated from all  $x_i$   
and their convex hull.

Randomization is essential.



$$\text{Vol}(K) = \text{Vol}(B) \cdot \mathbb{P}_2(x \in K)$$

DFK algorithm.  $K_i = K \cap 2^{i/n} B$

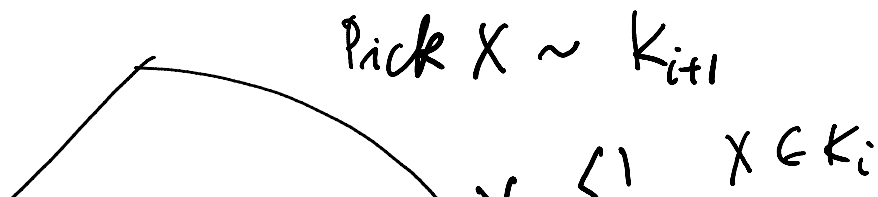
LV algorithm.

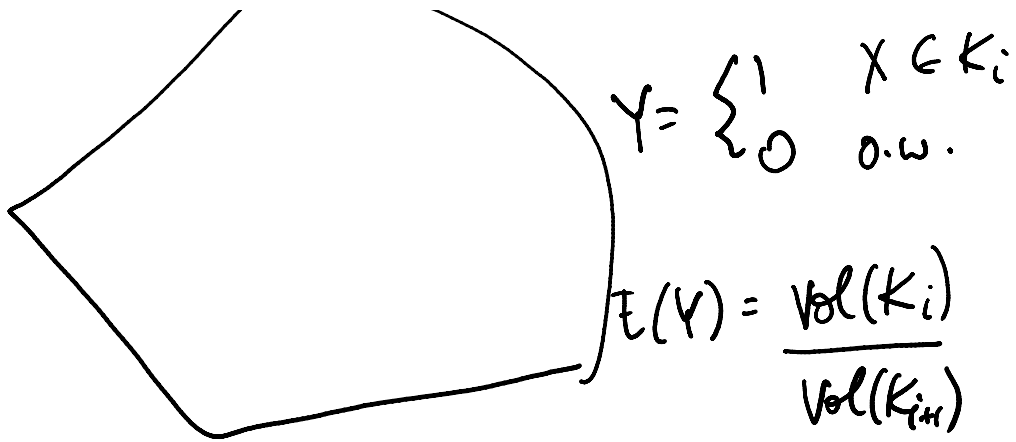
CV algorithm.

Volume / integration  $\longrightarrow$  Sampling.

$$\text{Vol}(K) = \text{Vol}(B_n) \cdot \frac{\text{Vol}(K_1)}{\text{Vol}(B_n)} \cdot \dots \cdot \frac{\text{Vol}(K_{i+1})}{\text{Vol}(K_i)} \dots$$

Estimate each ratio using random samples.





Lemma  $\text{Vol}(K_{i+1}) \leq 2 \text{Vol}(K_i)$

Pf.

$$\begin{aligned}
 K_{i+1} &= 2^{\frac{i+1}{n}} B_n \cap K \\
 &\subseteq 2^{\frac{1}{n}} (2^{i/n} B_n \cap K) \\
 &\subseteq 2^{\frac{1}{n}} K_i
 \end{aligned}$$

$$W_i = \frac{1}{k} \sum_{i=1}^k Y_i$$

$$V = \prod_{i=1}^m W_i$$

$$\begin{aligned}
 \text{Var}(V) &= \frac{E(V^2)}{E(V)^2} - 1 \\
 &= \prod_{i=1}^m \frac{E(W_i^2)}{E(W_i)^2} - 1
 \end{aligned}$$

$$= \prod_i \left( 1 + \frac{\text{Var}(W_i)}{E(W_i)^2} \right) - 1$$

Suffice to set  $\frac{\text{Var}(W_i)}{E(W_i)^2} \leq \frac{c \cdot \epsilon^2}{m}$

to get  $\text{Var}(V) \leq \epsilon^2$ .

$\Rightarrow K = O\left(\frac{m}{\epsilon^2}\right)$  in each phase

Total # samples =  $O\left(\frac{m^2}{\epsilon^2}\right) = O\left(\frac{n^2 \log^2 R}{\epsilon^2}\right)$

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How to sample?

1. Grid walk
  2. Ball walk
  3. Hit-and-Run
  4. Dikin, Geodesic, Hamiltonian (Polytopes)
- (general)

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Stationary distribution is uniform